Computing Cylindrical Algebraic Decomposition via Triangular Decomposition

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Background

Cylindrical algebraic decomposition (CAD) is a fundamental tool in real algebraic geometry. It was introduced by Collins in 1973 and has been followed by lots of improvements, like

- improved projection methods
  (McCallum 88, 98, Hong 90, Brown 01)
- partially built CADs
  (Collins and Hong 91, McCallum 93, Strzeboński 00)
- improved stack construction
  (Collins, Johnson and Krandick 02)
- efficient projection orders
  (Dolzmann, Seidl and Sturm 04)
- ...

Motivation

1. Understand the relations and possible interactions between CAD and triangular decompositions of polynomial systems.

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2. Investigate the possibility of improving the practical efficiency of CAD implementation by means of modular methods and fast polynomial arithmetic, being developed for triangular decompositions (Xin Li, Marc Moreno Maza and Wei Pan, ISSAC 09).

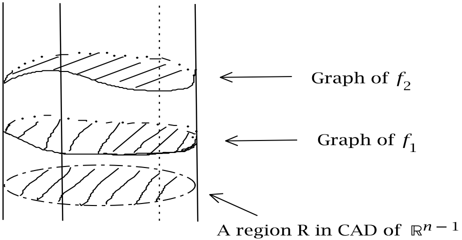
Cylindrical Algebraic Decomposition (I)

A cylindrical algebraic decomposition of $\mathbb{R}^n$ can be defined inductively as follows.

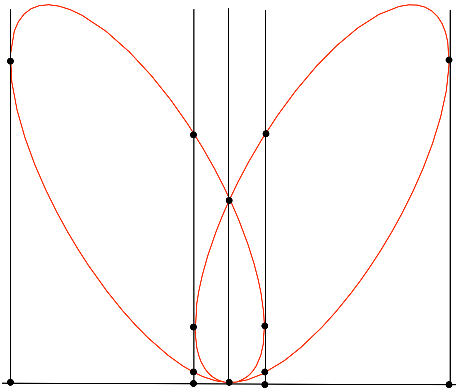
• $n = 1$. A CAD of $\mathbb{R}$ is a finite partition of the real line into points and open intervals.

Cylindrical Algebraic Decomposition (II)

- $n > 1$. Given a CAD D' of $\mathbb{R}^{n-1}$, one builds a CAD D of $\mathbb{R}^n$ as follows. Above each region R of D':
  - consider finitely many disjoint graphs (called *sections*) of continuous real-valued algebraic functions,
  - decomposing the cylinder $R \times \mathbb{R}^1$, into sections and *sectors* (located between two consecutive sections), which form a stack over R,
  - then all the sections and sectors are the elements of D.



A Cylindrical Algebraic Decomposition of $\mathbb{R}^2$ Induced by the Tacnode Curve



Tacnode curve: $y^4 - 2y^3 + y^2 - 3x^2y + 2x^4 = 0.$

Algorithm of Collins

Projection: Starting from the input $F_n \subset \mathbb{Q}[y_1, \ldots, y_n]$, repeatedly apply a projection operator to eliminate the variables one by one until a set of univariate polynomials are obtained

$$F_n \to F_{n-1} \to \cdots \to F_1$$

such that an F_k-invariant CAD of $\mathbb{R}^k$ can be constructed from an F_{k-1}-invariant CAD of $\mathbb{R}^{k-1}$, for $2 \leq k \leq n$.

Lifting: One isolates the real roots of polynomial in F_1 and deduces a CAD of $\mathbb{R}^1$. For each region of the CAD of $\mathbb{R}^1$, one evaluates the polynomials of F_2 at a *sample point* and isolates their real roots, from which one produces a stack over the region. Continuing in this manner, one finally obtains a CAD of $\mathbb{R}^n$.

Another View of CAD

A CAD of $\mathbb{R}^n$ is a partition of $\mathbb{R}^n$, where

- all the cells are cylindrically arranged, that is for all $1 \le j < n$ the projections on the first j coordinates $(y_1, \ldots, y_j)$ of any two cells are either identical or disjoint
- each cell is a connected semi-algebraic subset, called a region

For $F_n \subset \mathbb{Q}[y_1, \ldots, y_n]$, a CAD of $\mathbb{R}^n$ is F_n-invariant if in each region of it, the sign of each $f \in F_n$ is constant.

Our Method

F_n: a set of polynomials of $\mathbb{Q}[y_1, \ldots, y_n]$.

Initial Partition: we decompose $\mathbb{C}^n$ into disjoint constructible sets $C_1, \ldots, C_e$ such that for each $f \in F_n$, either f is identically zero in C_i or f vanishes at no points of C_i.

Make Cylindrical: we transform the initial partition and obtain another partition of $\mathbb{C}^n$ into disjoint constructible sets s.t. this second decomposition is cylindrical. We call it a cylindrical decomposition of $\mathbb{C}^n$.

Make Semi-Algebraic: from the previous decomposition we produce an F_n-invariant CAD of $\mathbb{R}^n$ via real root isolation of zero-dimensional regular chains.

The Three Phases

$F_n \subset \mathbb{Q}[y_1, \cdots, y_n]$

$\downarrow$

Initial Partition

$\downarrow$

$\mathcal{C}$: a partition of $\mathbb{C}^n$ into constructible sets

$\downarrow$

Make Cylindrical

$\downarrow$

$\mathcal{D}$: a cylindrically arranged partition of $\mathbb{C}^n$ into constructible sets

$\downarrow$

Make SemiAlgebraic

$\downarrow$

An F_n-invariant CAD of $\mathbb{R}^n$

Representation of Constructible Sets

A pair $R = [T, h]$ is called a regular system if T is a regular chain, and h is a polynomial which is regular w.r.t $\mathrm{sat}(T)$.

Theorem (CGLMP, CASC2007)

Every constructible set can be written as a finite union of the zero sets of regular systems.

The constructible set

$$\left\{ \begin{array}{rcl} x(1+y)-s & = & 0 \\ y(1+x)-s & = & 0 \\ x+y-1 & \neq & 0 \end{array} \right. \quad (1)$$

can be represented by two regular systems

$$R_1 : \left| \begin{array}{l} T_1 = \left\{ \begin{array}{l} (y+1)x - s \\ y^2 + y - s \end{array} \right. \\ h_1 = y - 2s + 1 \end{array} \right. \qquad R_2 : \left| \begin{array}{l} T_2 = \left\{ \begin{array}{l} x+1 \\ y+1 \\ s \end{array} \right. \\ h_2 = 1 \end{array} \right.$$

Initial Partition

- Decompose each hypersurface $f_i = 0$ as union of zero sets of regular systems.
- Compute an intersection free basis of the $s+1$ constructible sets $f_1 = 0, \ldots, f_s = 0$ and $f_1 \cdots f_s \neq 0$.

Consider the parametric parabola $p = ax^2 + bx + c$, where $x > c > b > a$. InitialPartition decomposes $\mathbb{C}^4$ into four pairwise disjoint sets, each of which is the zero set of a regular system.

$$r_1 := \left\{ \begin{array}{rcl} c & = & 0 \\ b & = & 0 \\ a & = & 0 \end{array} \right. , \quad r_2 := \left\{ \begin{array}{rcl} bx + c & = & 0 \\ b & \neq & 0 \\ a & = & 0 \end{array} \right. ,$$

$$r_3 := \left\{ \begin{array}{rcl} ax^2 + bx + c & = & 0 \\ a & \neq & 0 \end{array} \right. , \quad r_4 := \left\{ \begin{array}{l} ax^2 + bx + c \neq 0 \end{array} \right. .$$

The Three Phases

$$F_n \subset \mathbb{Q}[y_1, \cdots, y_n]$$

$\downarrow$

Initial Partition

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Make SemiAlgebraic

$\downarrow$

An F_n-invariant CAD of $\mathbb{R}^n$

Separate Zeros

Let $rs = [T, h]$ be a regular system of $\mathbb{Q}[y_1 < \cdots < y_n]$. We see $y_1, \ldots, y_{n-1}$ as parameters, denoted by $\mathbf{u}$, and regard rs as a parametric system in $\mathbf{u}$, that we solve via comprehensive triangular decomposition (CGLMP 07).

It computes a partition of the projection onto the $\mathbf{u}$-space of $Zero(rs)$ such that, above each cell R of the partition, $Zero(rs)$ equals the union of the zero sets of some polynomials $p_1, \ldots, p_r \in \mathbb{R}[y_1, \ldots, y_n]$, where

- the initial of each p_j does not vanish on R,
- the p_j's are squarefree and pairwise coprime at any point of R.

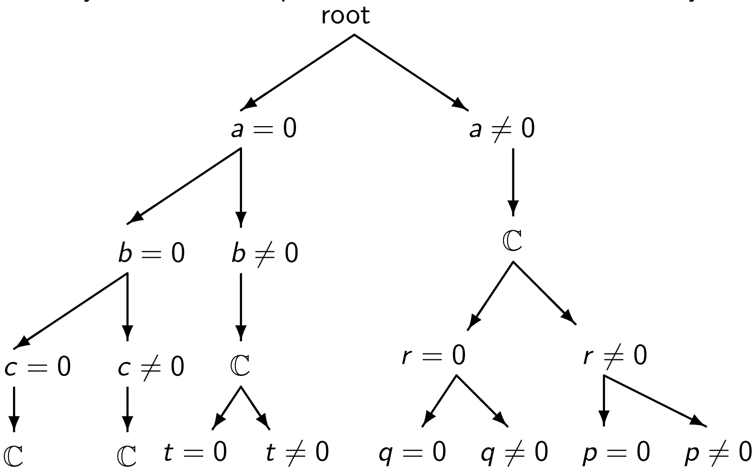
For the regular system

$$r_3 := \begin{cases} ax^2 + bx + c & = & 0 \\ a & \neq & 0 \end{cases}$$

Calling SeparateZeros(r_3) will get

$$\begin{array}{lcc} \text{polynomials :} & ax^2 + bx + c & 2ax + b \\ & \uparrow & \uparrow \\ \mathbf{u}-\text{space :} & a(4ac - b^2) \neq 0 & 4ac - b^2 = 0 \ \wedge \ a \neq 0 \end{array}$$

Let $t = bx + c$, $q = 2ax + b$, $r = 4ac - b^2$ and $p = ax^2 + bx + c$. A p-invariant cylindrical decomposition of $\mathbb{C}^4$ can be described by a tree.



The Three Phases

$$F_n \subset \mathbb{Q}[y_1, \cdots, y_n]$$

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Make SemiAlgebraic

$$\downarrow$$

An F_n-invariant CAD of $\mathbb{R}^n$

Make SemiAlgebraic (I)

Theorem (Collins)

Let $p \in \mathbb{R}[y_1, \ldots, y_n]$ and R be a region of $\mathbb{R}^{n-1}$. If init(p) does not vanish on R and the number of distinct complex roots of p is invariant on R, then p is delineable *on R, that is, $V(p)$ is the union of finitely many disjoint graphs of continuous functions over R.*

Make SemiAlgebraic (II)

Corollary

Let $\{p_1, \ldots, p_r\} \subset \mathbb{Q}[y_1 < \cdots < y_n]$ *and let* R *be a region of* $\mathbb{R}^{n-1}$*. Assume that for all* $\alpha \in R$*:*

- *the initial of each* p_i *does not vanish at* α*;*

Make SemiAlgebraic (II)

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- *all* $p_i(\alpha, y_n)$ *are squarefree and pairwise coprime.*

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Then each p_i *is delineable on* R *and any two sections of the cylinder over* R*, given by different* p_i *and* p_j*, are disjoint.*

Make SemiAlgebraic (II)

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By Collins' theorem and its corollary, one derives a CAD of $\mathbb{R}^n$ from a cylindrical decomposition of $\mathbb{C}^n$, by means of real root isolation of zero-dimensional regular chains (Bican Xia and Ting Zhang 06).

Comparing with Collins' Algorithm

Consider the parametric parabola $p = ax^2 + bx + c$, where $x > c > b > a$.

- Our algorithm produces a p-invariant CAD of $\mathbb{R}^4$ with 27 cells, which is minimal.
- By Collins-Hong or McCallum projection operator, one produces the following polynomials during the projection phase:

  $$ax^2 + bx + c, b^2 - 4ac, c, b, a.$$

  In the lifting phase, one then obtains a CAD of $\mathbb{R}^4$ with 115 cells (Brown 01)!
- If Brown-McCallum projection operator is applied, one could also obtain a CAD of $\mathbb{R}^4$ with 27 cells (Brown 01). However, this projection operator may fail in some (rare) cases.

Sys	InitialPartition	MakeCylindrical	MakeSemiAlgebraic	Total	$N_{\mathbb{R}}$
1	0.024	0.096	0.024	0.144	27
2	1.184	2.856	1.048	5.088	895
3	0.004	7.512	0.704	8.220	233
4	0.264	1.368	1.080	2.716	421
5	0.016	0.052	0.116	0.184	55
6	0.108	0.156	0.120	0.384	41
7	2.704	3.600	1.360	7.664	893
8	0.380	1.608	1.196	3.184	365
9	0.288	0.532	0.264	1.084	209
10	5.668	48.079	18.833	72.640	3677
11	0.252	1.192	0.620	2.068	563
12	2.664	135.028	88.142	225.862	20143
13	10.576	35.846	6.905	53.335	4949
14	5.728	71.760	2520.354	2597.878	27547
15	690.731	2513.817	299.250	3503.954	66675
16	895.435	2064.469	-	-	-
17	0.052	-	-	-	-
18	-	-	-	-	-

Table 1 Timing (s) and number of cells for CAD

Observation

- For most examples the steps of the algorithm dedicated to computations in $\mathbb{C}^n$, where GCDs of polynomials modulo regular chains are computed intensively, dominate the step taking place in $\mathbb{R}^n$.
- The data suggests that the modular methods and efficient implementation techniques being developed in `RegularChains` library have a large potential for improving our current implementation.

Conclusion and Future Work

- We have presented a new algorithm for computing CAD based on triangular decomposition techniques,
- We have introduced an intermediate concept, cylindrical decomposition of the complex space, from which a CAD of $\mathbb{R}^n$ can easily be extracted.

Some future work:

- Improve our black-box subroutine in the algorithm level,
- Integrate fast arithmetic and modular methods into our algorithm,
- Quantifier elimination.

1. Parametric parabola: $\{ax^2 + bx + c\}, x > c > b > a$.
2. Whitney umbrella: $\{x - uv, y - v, z - u^2\}, v > u > z > y > x$.
3. Quartic: $\{x^4 + px^2 + qx + r\}, x > p > q > r$.
4. Sphere-Catastrophe: $\{z^2 + y^2 + x^2 - 1, z^3 + xz + y\}, x > y > z$.
5. Tacnode curve: $\{y^4 - 2y^3 + y^2 - 3x^2y + 2x^4\}, y > x$.
6. Arnon-84-2: $\{144y^2 + 96x^2y + 9x^4 + 105x^2 + 70x - 98,$
$xy^2 + 6xy + x^3 + 9x\}, y > x$.
7. A real implicitization problem:
$\{x - uv, y - uv^2, z - u^2\}$, $v > u > z > y > x$.
8. Ball-circular-cylinder:
$\{x^2 + y^2 + z^2 - 1, x^2 + (y + z - 2)^2 - 1\}, z > y > x$.
9. Termination of term rewrite system
$\{x - r, y - r, x^2(1 + 2y)^2 - y^2(1 + 2x^2)\}, r > x > y$.
10. Collins and Johnson: $\{3a^2r + 3b^2 - 2ar - a^2 - b^2,$
$3a^2r + 3b^2r - 4ar + r - 2a^2 - 2b^2 + 2a,\ a - 1/2, b, r, r - 1\},$
$r > a > b$.
11. Range of lower bounds
$\{a, az^2 + bz + c, ax^2 + bx + c - y\}$, $z > c > b > a > x > y$.

12. X-axis ellipse problem: $\{b^2(x-c)^2+a^2y^2-a^2b^2,$
$x^2+y^2-1\}, y>x>b>c>a.$

13. Davenport and Heintz
$\{a-d, b-c, a-c, b-1, a^2-b\}, a>b>c>d.$

14. Hong-90
$\{r+s+t, rs+st+tr-a, rst-b\},\ t>s>r>b>a.$

15. Solotareff-3
$\{r, r-1, u+1, u-v, v-1,\ 3u^2+2ru-a, 3v^2+2rv-a,$
$u^3+ru^2-au+a-r-1,\ v^3+rv^2-av-2b-a+r+1\},$
$b>u>v>r>a.$

16. Collision problem
$\{\frac{17}{16}t-6, \frac{17}{16}t-10, x-\frac{17}{16}t+1,\ x-\frac{17}{16}t-1,\ y-\frac{17}{16}t+9,$
$y-\frac{17}{16}t+7,\ (x-t)^2+y^2-1\}, t>x>y.$

17. McCallum trivariate random polynomial
$\{(y-1)z^4+xz^3+x(1-y)z^2+(y-x-1)z+y\},\ z>y>x.$

18. Ellipse problem
$\{b^2(x-c)^2+a^2(y-d)^2-a^2b^2, a, b, x^2+y^2-1\},$
$y>x>d>c>b>a.$